

# FIRST EXPERIMENTS WITH LEARNING INSTANTIATIONS

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Konstantin Korovin    Josef Urban

University of Manchester

Czech Technical University in Prague

ICMS 2018  
July 27, 2018



# Motivation: Lemmatization and Instantiation

- Lemmatization:
  - Add to a problem lemmas useful in previous proofs
- Instantiation:
  - Add to a problem instances useful in previous proofs

# Previous Lemmatization Experiments - HOL Light and Flyspeck

- Over 1B low-level (proof) lemmas in Flyspeck
- 1.5M-7M higher-level lemmas in Flyspeck
- Define fast preprocessing methods to extract the most important ones:
- PageRank, recursive dependency count, recursive use count, etc.
- Use the most important lemmas together with the toplevel theorems
- helps by 5-20% on HOL Light/Flyspeck
- Quite often the lemmas are instances or easy consequences

# Example of Useful Flyspeck Proof Lemmas

Conjecture:

```
AFFINE_ALT: |- affine s <=> (!x y u. x IN s /\ y IN s  
    ==> (&1 - u) % x + u % y IN s)
```

Easy, but useful lemmas (derived inside previous proofs):

```
NEWDEP309638: |- &1 - a + a = &1
```

```
NEWDEP310357: |- -- &1 * -- &1 = &1
```

```
NEWDEP272099_conjunct1: |- !m. &m + -- &m = &0
```

## Previous Lemmatization Experiments - E prover

- refutational TPTP proofs of many Mizar problems
- postprocess (redirect) them to obtain direct proofs
- collect the proof lemmas derived only from the axioms
- Use them along with the toplevel premises for proving new conjectures
- 6% improvement on a Mizar-based benchmark

# Example of Useful Mizar Proof Lemmas

Lemma:

$$X1 \setminus (X2 \setminus X1) = X1$$

proved in

$$X \subseteq Y \text{ implies } Z \setminus Y \subseteq Z \setminus X$$

useful in

$$X \setminus (Y \setminus X) = X \setminus Y$$

$$(X \setminus Y) \setminus Z = (X \setminus Z) \setminus (Y \setminus Z)$$

etc.

# Why Learn Instantiation?

- Many of the proof lemmas in the previous experiments are useful instances
- A targeted instance prevents redundant inferences
- Instantiation-based ATP calculi are reaching state of the art
- SMT: quantifier instantiation
- iProver: calculus based on gradual instantiation (and SAT)

- instantiate (abstract) all clauses with a unique constant and call a SAT solver
- if unsatisfiable, we are done
- else add more complicated ground instances – the Inst-Gen rule:
- from  $L \vee C, \neg L' \vee D$  create  $(L \vee C)\theta, (\neg L' \vee D)\theta$
- and continue
- Creation of the new instances is guided by the propositional assignment.
- In particular, we instantiate clauses with unifiable literals that have different value in the ground model
- However, this is often still very non-deterministic



# Our Idea and Plan

- If we immediately guess the right instantiations, we are done (if it's SAT-easy)
- How do we guess them?
- Machine learning from many related problems!
- Our research plan:
  - **Add all previously useful clause instances**
  - Train a ranker and add only some number of most useful clause instances for the current context (analogous to premise selection)
  - Generate new instances by learning useful terms for the current context (another variant of premise selection)
  - Generate new terms for the current context by a trained probabilistic grammar

# First Experiment

- Run standard iProver on 32525 Mizar problems
- solves 21685 problems with 10s time limit
- from the 21685 proofs extract the useful instances for each initial clause  $C$
- whenever an unsolved problem contains  $C$ , add *all* the previously useful instances
- rerun iProver on the unsolved problems
- This adds 621 problems
- Many problems are too big - we really need the pruning

# Future and Thanks

- That's all we did for now
- Encouraging, but we need to proceed with the next steps
- Thanks for your attention!