DEVELOPMENTS IN AI AND THEOREM PROVING

Josef Urban

Czech Technical University in Prague

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How Do We Automate Math and Science?

- What is mathematical and scientific thinking?
- · Pattern-matching, analogy, induction from examples
- · Deductive reasoning
- · Complicated feedback loops between induction and deduction
- Using a lot of previous knowledge both for induction and deduction
- · We need to develop such methods on computers
- · Are there any large corpora suitable for nontrivial deduction?
- · Yes! Large libraries of formal proofs and theories
- So let's develop strong AI on them!

Why Combine Learning and Reasoning Today?

1 It practically helps!

- Automated theorem proving for large formal verification is useful:
 - Formal Proof of the Kepler Conjecture (2014 Hales 20k lemmas)
 - · Formal Proof of the Feit-Thompson Theorem (2012 Gonthier)
 - · Verification of compilers (CompCert) and microkernels (seL4)
 - · Verification hardware architectures, transport systems, trading rules
 - ...
- · But good learning/AI methods needed to cope with large theories!

2 Blue Sky Al Visions:

- Get strong AI by learning/reasoning over large KBs of human thought?
- · Big formal theories: good semantic approximation of such thinking KBs?
- · Deep non-contradictory semantics better than scanning books?
- Gradually try learning math/science:
 - · What are the components (inductive/deductive thinking)?
 - · How to combine them together?
 - · Automate/verify math, science, law, ...
 - · Leibniz: Calculemus resolve disputes
 - J. McCarthy: Mathematical Objectivity and the Power of Initiative

What is Formal Mathematics?

- · Developed thanks to the Leibniz/Russell/Frege/Hilbert/... program
- Mathematics put on formal logic foundations (symbolic computation)
- ... which btw. led also to the rise of computers (Turing/Church, 1930s)
- Formal math (1950/60s): combine formal foundations and the newly available computers
- De Bruijn, Milner, Trybulec, Boyer and Moore, Gordon, Huet, Paulson, ...
- · Automath, LCF, Mizar, NQTHM and ACL2, HOL, Coq, Isabelle, ...
- Conceptually very simple:
- · Write all your axioms and theorems so that computer understands them
- Write all your inference rules so that computer understands them
- · Use the computer to check that your proofs follow the rules
- But in practice, it turns out not to be so simple
- · Many approaches, still not mainstream, but big breakthroughs recently

Freek Wiedijk's Example: Irrationality of $\sqrt{2}$ (informal text)

tiny proof from Hardy & Wright:

Theorem 43 (Pythagoras' theorem). $\sqrt{2}$ is irrational. The traditional proof ascribed to Pythagoras runs as follows. If $\sqrt{2}$ is rational, then the equation

$$a^2 = 2b^2$$
 (4.3.1)

is soluble in integers *a*, *b* with (a, b) = 1. Hence a^2 is even, and therefore *a* is even. If a = 2c, then $4c^2 = 2b^2$, $2c^2 = b^2$, and *b* is also even, contrary to the hypothesis that (a, b) = 1.

Irrationality of $\sqrt{2}$ (Formal Proof Sketch)

exactly the same text in Mizar syntax:

```
theorem Th43: :: Pythagoras' theorem
  sqrt 2 is irrational
proof
  assume sqrt 2 is rational;
  consider a,b such that
4 3 1: a^2 = 2 \cdot b^2 and
    a,b are relative prime;
  a^2 is even;
  a is even;
  consider c such that a = 2 * c;
  4 \star c^2 = 2 \star b^2;
  2 \star c^2 = b^2;
  b is even;
  thus contradiction;
end;
```

Irrationality of $\sqrt{2}$ in HOL Light

let SQRT_2_IRRATIONAL = prove (`~rational(sqrt(&2))`, SIMP_TAC[rational; real_abs; SQRT_POS_LE; REAL_POS] THEN REWRITE_TAC[NOT_EXISTS_THM] THEN REPEAT GEN_TAC THEN DISCH_THEN(CONJUNCTS_THEN2 ASSUME_TAC MP_TAC) THEN SUBGOAL_THEN `~((&p / &q) pow 2 = sqrt(&2) pow 2)` (fun th -> MESON_TAC[th]) THEN SIMP_TAC[SQRT_POW_2; REAL_OS; REAL_POW_DIV] THEN ASM_SIMP_TAC[REAL_EQ_LDIV_EQ; REAL_OF_NUM_LI; REAL_POW_LT; ARITH_RULE `0 < q <=> ~(q = 0)`] THEN ASM_MESON_TAC[NSQRT_2; REAL_OF_NUM_POW; REAL_OF_NUM_MUL; REAL_OF_NUM_EQ]);;

Irrationality of $\sqrt{2}$ in Isabelle/HOL

```
theorem sgrt2 not rational:
  "sort (real 2) ∉ 0"
proof
 assume "sqrt (real 2) \in \mathbb{Q}"
  then obtain m n :: nat where
    n_nonzero: "n \neq 0" and sqrt_rat: "!sqrt (real 2)! = real m / real n"
    and lowest_terms: "gcd m n = 1" ...
 from n_nonzero and sqrt_rat have "real m = {sqrt (real 2)} * real n" by simp
  then have "real (m^2) = (sort (real 2))^2 * real <math>(n^2)"
    by (auto simp add: power2 eg square)
  also have "(sgrt (real 2))<sup>2</sup> = real 2" by simp
  also have "... * real (m^2) = real (2 * n^2)" by simp
  finally have eq: m^2 = 2 * n^2 ...
  hence "2 dvd m<sup>2</sup>"...
  with two is prime have dvd m: "2 dvd m" by (rule prime dvd power two)
  then obtain k where "m = 2^* k"
  with eq have "2 * n^2 = 2^2 * k^2" by (auto simp add: power2 eq square mult ac)
  hence "n^2 = 2 * k^2" by simp
  hence "2 dvd n^2"...
  with two_is_prime have "2 dvd n" by (rule prime_dvd_power_two)
  with dvd m have "2 dvd qcd m n" by (rule qcd greatest)
  with lowest terms have "2 dvd 1" by simp
 thus False by arith
ged
```

Big Example: The Flyspeck project

• Kepler conjecture (1611): The most compact way of stacking balls of the same size in space is a pyramid.



- Formal proof finished in 2014
- · 20000 lemmas in geometry, analysis, graph theory
- All of it at https://code.google.com/p/flyspeck/
- · All of it computer-understandable and verified in HOL Light:
- polyhedron s /\ c face_of s ==> polyhedron c
- However, this took 20 30 person-years!

What Are Automated Theorem Provers?

- · Computer programs that (try to) determine if
 - A conjecture C is a logical consequence of a set of axioms Ax
 - · The derivation of conclusions that follow from facts by inference rules
- Systems: Vampire, E, SPASS, Prover9, Z3, CVC4, Satallax, ...
- Brute-force search calculi (resolution, superposition, tableaux, SMT, ...)
- · Human-designed heuristics for pruning of the search space
- Fast combinatorial explosion on large knowledge bases like Flyspeck and Mizar
- · Need to be equipped with good domain-specific inference guidance ...
- ... and that is what I try to do ...
- ... typically by learning in various ways from the knowledge bases ...

History and Motivation for AI/TP

- Intuition vs Formal Reasoning Poincaré vs Hilbert, Science & Method
- Turing's 1950 paper: Learning Machines, learn Chess?, undecidability??
- · Lenat, Langley, etc: manually-written heuristics, learn Kepler laws,...
- Denzinger, Schulz, Goller, Fuchs late 90's, ATP-focused:
- Learning from Previous Proof Experience
- · My MSc (1998): Try ILP to learn rules and heuristics from IMPS/Mizar
- · Since: Use large formal math (Big Proof) corpora: Mizar, Isabelle, HOL
- ... to combine/develop symbolic/statistical deductive/inductive ML/TP/AI
- ... hammer-style methods, feedback loops, internal guidance, ...
- More details AGI'18 keynote: https://bit.ly/3qifhg4
- Al vs DL: Ben Goertzel's Prague talk: https://youtu.be/Zt2HSTuGBn8
- Big Al visions: automate/verify math, science, law, (Leibniz, McCarthy, ..)
- Practical impact: boost today's large ITP verification projects

Using Learning to Guide Theorem Proving

- · high-level: pre-select lemmas from a large library, give them to ATPs
- · high-level: pre-select a good ATP strategy/portfolio for a problem
- high-level: pre-select good hints for a problem, use them to guide ATPs
- Iow-level: guide every inference step of ATPs (tableau, superposition)
- · low-level: guide every kernel step of LCF-style ITPs
- mid-level: guide application of tactics in ITPs
- mid-level: invent suitable ATP strategies for classes of problems
- mid-level: invent suitable conjectures for a problem
- mid-level: invent suitable concepts/models for problems/theories
- · proof sketches: explore stronger/related theories to get proof ideas
- theory exploration: develop interesting theories by conjecturing/proving
- feedback loops: (dis)prove, learn from it, (dis)prove more, learn more, ...
- autoformalization: (semi-)automate translation from LATEX to formal

Large AI/TP Datasets

- Mizar / MML / MPTP since 2003
- MPTP Challenge (2006), MPTP2078 (2011), Mizar40 (2013)
- Isabelle (and AFP) since 2005
- Flyspeck (including core HOL Light and Multivariate) since 2012
- HOL4 since 2014, CakeML 2017, GRUNGE 2019
- Coq since 2013/2016
- · AIM Veroff & Kinyon, Loops with Abelian Inner Mappings long proofs
- · Lean?, Stacks?, Arxiv?, ProofWiki?, ...

- ENIGMA/hammer proofs of Pythagoras: https://bit.ly/2MVPAn7 (more at http://grid01.ciirc.cvut.cz/~mptp/enigma-ex.pdf) and simplified Carmichael https://bit.ly/3oGBdRz
- Hammering demo: http://grid01.ciirc.cvut.cz/~mptp/out4.ogv
- TacticToe on HOL4:

http://grid01.ciirc.cvut.cz/~mptp/tactictoe_demo.ogv

Tactician for Coq:

https://blaauwbroek.eu/papers/cicm2020/demo.mp4, https://coq-tactician.github.io/demo.html

• Inf2formal over HOL Light:

http://grid01.ciirc.cvut.cz/~mptp/demo.ogv

High-level ATP guidance: Premise Selection

- Early 2003: Can existing ATPs be used over the freshly translated Mizar library?
- About 80000 nontrivial math facts at that time impossible to use them all
- · Is good premise selection for proving a new conjecture possible at all?
- Or is it a mysterious power of mathematicians? (Penrose)
- Today: Premise selection is not a mysterious property of mathematicians!
- · Reasonably good algorithms started to appear (more below).
- Extensive human (math) knowledge obsolete?? (cf. Watson, Debater, ..)
- Since 2004 (my PhD): many examples of nontrivial alternative proofs proposed by the Als in Mizar, Flyspeck, Isabelle, ...
- The premise selection algorithms see wider than humans

Today's AI-ATP systems (*-Hammers)



How much can it do?

- Mizar / MML MizAR
- Isabelle (Auth, Jinja) Sledgehammer
- Flyspeck (including core HOL Light and Multivariate) HOL(y)Hammer
- HOL4 (Gauthier and Kaliszyk), Coq (Czajka and Kaliszyk)
- · Rigorous resource controlled train/test evaluations on toplevel lemmas:

 \approx 40-45% success rate by 2016 \approx 60% on Mizar as of 2021

Premise Selection and Hammer Methods

- · Many syntactic features (symbols, walks in the parse trees)
- More semantic features encoding
- · Term matching/unification, validity in models, latent semantics (LSI)
- Distance-weighted k-nearest neighbor, SVMs, Naive Bayes
- · Gradient boosted decision trees (GBDTs XGBoost, LightGBM)
- · Neural models: CNNs, RNNs/Attention/Transformers/GPT, GNNs
- As of 2020, tough competition between GBDTs, GNNs and RNNs/Transformers (and relatives)
- · K-NN still very good, Olsak's logic-aware GNN probably best
- RNNs/Transformers good at stateful premise selection (Piotrowski 2019,2020)
- · Ensemble methods combining the different predictors help a lot

Premise Selection and Hammer Methods

- · Learning in a binary setting from many alternative proofs
- Interleaving many learning and proving runs (*MaLARea loop 2006*) to get positives/negatives (ATPBoost - Piotrowski 2018)
- Matching and transferring concepts and theorems between libraries (Gauthier & Kaliszyk) allows "superhammers", conjecturing, and more
- Lemmatization extracting and considering millions of low-level lemmas and learning from their proofs (Kaliszyk & JU 2013)
- Hammers combined with guided tactical search: TacticToe (Gauthier HOL4) and its later relatives

FACE_OF_POLYHEDRON_POLYHEDRON

```
let FACE OF POLYHEDRON POLYHEDRON = prove
 ('!s:real^N->bool c. polyhedron s /\ c face of s ==> polyhedron c',
 REPEAT STRIP TAC THEN FIRST ASSUM
   (MP TAC O GEN REWRITE RULE I [POLYHEDRON INTER AFFINE MINIMAL]) THEN
  REWRITE TAC[RIGHT IMP EXISTS THM; SKOLEM THM] THEN
  SIMP TAC[LEFT IMP EXISTS THM; RIGHT AND EXISTS THM; LEFT AND EXISTS THM] THEN
 MAP EVERY X GEN TAC
   ['f:(real^N->bool)->bool'; 'a:(real^N->bool)->real^N';
    'b: (real^N->bool) ->real'] THEN
  STRIP TAC THEN
 MP TAC(ISPECL ['s:real^N->bool': 'f:(real^N->bool)->bool':
                 `a:(real^N->bool)->real^N`; `b:(real^N->bool)->real`]
         FACE OF POLYHEDRON EXPLICIT) THEN
 ANTS TAC THENL [ASM REWRITE TAC]] THEN ASM MESON TAC]]; ALL TAC] THEN
  DISCH THEN (MP TAC o SPEC 'c:real^N->bool') THEN ASM REWRITE TAC[] THEN
 ASM CASES TAC 'c:real^N->bool = {}' THEN
 ASM REWRITE TAC[POLYHEDRON EMPTY] THEN
 ASM CASES TAC 'c:real^N->bool = s' THEN ASM REWRITE TAC[] THEN
  DISCH THEN SUBST1 TAC THEN MATCH MP TAC POLYHEDRON INTERS THEN
  REWRITE TAC[FORALL IN GSPEC] THEN
 ONCE REWRITE TAC[SIMPLE IMAGE GEN] THEN
 ASM SIMP TAC[FINITE IMAGE: FINITE RESTRICT] THEN
 REPEAT STRIP TAC THEN REWRITE TAC[IMAGE ID] THEN
 MATCH MP TAC POLYHEDRON INTER THEN
 ASM REWRITE TAC[POLYHEDRON HYPERPLANE]);;
```

FACE_OF_POLYHEDRON_POLYHEDRON

polyhedron s /\ c face_of s ==> polyhedron c

HOL Light proof: could not be re-played by ATPs.

Alternative proof found by a hammer based on FACE_OF_STILLCONVEX: Face *t* of a convex set *s* is equal to the intersection of *s* with the affine hull of *t*.

```
FACE_OF_STILLCONVEX:
 !s t:real^N->bool. convex s ==>
 (t face_of s <=>
  t SUBSET s /\ convex(s DIFF t) /\ t = (affine hull t) INTER s)
POLYHEDRON_IMP_CONVEX:
 !s:real^N->bool. polyhedron s ==> convex s
POLYHEDRON_INTER:
 !s t:real^N->bool. polyhedron s /\ polyhedron t
 ==> polyhedron (s INTER t)
POLYHEDRON_AFFINE_HULL:
 !s. polyhedron(affine hull s)
```

Statistical Guidance of a Simple Connection Prover

- · learn guidance of every clausal inference in connection tableau (leanCoP)
- · set of first-order clauses, extension and reduction steps
- · proof finished when all branches are closed
- · a lot of nondeterminism, requires backtracking
- the search space quickly explodes
- · good for learning the tableau compactly represents the proof state



Using Reinforcement Learning to Guide leanCoP

- Monte-Carlo Tree Search (MCTS) used in AlphaGo
- MCTS search nodes are sequences of clause application
- a good heuristic to explore new vs exploit good nodes:

$$\frac{w_i}{n_i} + c \cdot p_i \cdot \sqrt{\frac{\ln N}{n_i}}$$
 (UCT - Kocsis, Szepesvari 2006)

- we learn the *policy* clause selection
- ... and the value proof state evaluation
- · big issue: representing clauses and proofs for learning
- many approaches none too good yet, esp. for value
- deep learning not impressive yet and slower than GBDTs
- · feedback loop between proving and learning many iterations

Statistical Guidance of Connection Tableau - rlCoP

- On 32k Mizar40 problems using 200k inference limit
- nonlearning CoPs:

| System | leanCoP | bare prover | rlCoP no policy/value (UCT only) |
|--------------------------|-------------|-------------|----------------------------------|
| Training problems proved | 10438 | 4184 | 7348 |
| Testing problems proved | 1143 | 431 | 804 |
| Total problems proved | 11581 | 4615 | 8152 |

- · rICoP with policy/value after 5 proving/learning iters on the training data
- 1624/1143 = 42.1% improvement over leanCoP on the testing problems

| Iteration | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|-----------------|-------|-------|-------|-------|-------------|-------|-------|--------------|
| Training proved | 12325 | 13749 | 14155 | 14363 | 14403 | 14431 | 14342 | 14498 |
| Testing proved | 1354 | 1519 | 1566 | 1595 | 1624 | 1586 | 1582 | 1591 |

Recent CoP Mutants: FLoP, GNN, RNN, lazyCoP

FLoP – Finding Longer Proofs (Zombori et al, 2019)



- Curriculum Learning used for connection tableau over Robinson Arithmetic
- addition and multiplication learned perfectly from 1 * 1 = 1
- · headed towards learning algorithms/decision procedures from math data
- currently black-box, combinations with symbolic methods (ILP) our next target
- Using RNNs for better tableau encoding, prediction of actions ...
- ... even guessing (decoding) next tableau literals (Piotrowski 2020)
- plCoP (Zombori 20), GNN-CoP (Olsak 20), lazyCoP (Rawson)
- · Zombori: learning new explainable Prolog actions (tactics) from proofs

ENIGMA: Guiding the Best ATPs like E Prover

- · Similar to rICoP interleave proving and learning of ENIGMA guidance
- · resolution/superposition harder for learning than tableau
- the proof state are two large heaps of clauses processed/unprocessed
- Done on 57880 Mizar problems recently 6 prove/learn iterations
- · Feedback loop: 70% improvement over the original strategy in 2019
- From 14933 proofs to 25397 proofs (all 10s CPU no cheating)
- Went up to 40k proofs in more iterations and 60s time in 2020
- Many proof examples at https://github.com/ai4reason/ATP_Proofs

| | S | $\mathcal{S} \odot \mathcal{M}_9^0$ | $\mathcal{S} \oplus \mathcal{M}_9^0$ | $\mathcal{S} \odot \mathcal{M}_9^1$ | $\mathcal{S} \oplus \mathcal{M}_9^1$ | $\mathcal{S} \odot \mathcal{M}_9^2$ | $\mathcal{S} \oplus \mathcal{M}_9^2$ | $S \odot \mathcal{M}_9^3$ | $\mathcal{S} \oplus \mathcal{M}_9^3$ |
|----------------|-------|-------------------------------------|--------------------------------------|-------------------------------------|--------------------------------------|-------------------------------------|--------------------------------------|---------------------------|--------------------------------------|
| solved | 14933 | 16574 | 20366 | 21564 | 22839 | 22413 | 23467 | 22910 | 23753 |
| S% | +0% | +10.5% | +35.8% | +43.8% | +52.3% | +49.4% | +56.5% | +52.8% | +58.4 |
| $\mathcal{S}+$ | +0 | +4364 | +6215 | +7774 | +8414 | +8407 | +8964 | +8822 | +9274 |
| $\mathcal{S}-$ | -0 | -2723 | -782 | -1143 | -508 | -927 | -430 | -845 | -454 |

| | $S \odot \mathcal{M}_{12}^3$ | $\mathcal{S} \oplus \mathcal{M}^3_{12}$ | $S \odot \mathcal{M}^3_{16}$ | $S \oplus \mathcal{M}^3_{16}$ |
|-----------------|------------------------------|---|------------------------------|-------------------------------|
| solved | 24159 | 24701 | 25100 | 25397 |
| $\mathcal{S}\%$ | +61.1% | +64.8% | +68.0% | +70.0% |
| $\mathcal{S}+$ | +9761 | +10063 | +10476 | +10647 |
| S- | -535 | -295 | -309 | -183 |

ENIGMA Proof Example – Knaster fixed-point theorem

```
theorem Th21:
 ex a st a is a fixpoint of f
  set H = {h where h is Element of L: h [= f.h};
  set fH = {f.h where h is Element of L: h [= f.h};
  set uH = "\/"(H, L);
 set fuH = "\/"(fH, L);
 take uH;
  now
   let fh be Element of L;
   assume fh in fH;
   then consider h being Element of L such that
Al: fh = f.h and
A2: h [= f.h;
   h in H by A2;
   then h [= uH by LATTICE3:38;
   hence fh [= f.uH by Al,QUANTAL1:def 12;
  end;
  then fH is_less_than f.uH by LATTICE3:def 17;
  then
A3: fuH [= f.uH by LATTICE3:def 21;
  now
    let a be Element of L:
   assume a in H;
    then consider h being Element of L such that
A4: a = h \& h [= f.h;
    reconsider fh = f.h as Element of L:
    take fh;
    thus a [= fh & fh in fH by A4;
  end;
  then uH [= fuH by LATTICE3:47;
  then
A5: uH [= f.uH by A3, LATTICES: 7;
  then f.uH [= f.(f.uH) by QUANTAL1:def 12;
  then f.uH in H;
  then f.uH [= uH by LATTICE3:38;
 hence uH = f.uH by A5, LATTICES:8;
end;
```

Low-level Symbolic ATP guidance: Prover9 hints

- The Prover9 community: non-associative algebra, 20-100k long proofs
- · Hints (Bob Veroff): extract lemmas from easier proofs to guide new proofs
- The hints behave like checkpoints you are on the right track
- · Gluing the different ideas from many other proofs together by search
- Exploration to get good hints (not really automated yet)
- · Very recent huge breakthrough in the AIM project by Kinyon and Veroff

TacticToe: mid-level ITP Guidance (Gauthier'17,18)

- TTT learns from human and its own tactical HOL4 proofs
- · No translation or reconstruction needed native tactical proofs
- · Fully integrated with HOL4 and easy to use
- · Similar to rICoP: policy/value learning for applying tactics in a state
- · However much more technically challenging a real breakthrough:
 - · tactic and goal state recording
 - · tactic argument abstraction
 - absolutization of tactic names
 - nontrivial evaluation issues
 - · these issues have often more impact than adding better learners
- · policy: which tactic/parameters to choose for a current goal?
- · value: how likely is this proof state succeed?
- 66% of HOL4 toplevel proofs in 60s (better than a hammer!)
- · similar work for Isabelle (Nagashima 2018), HOL Light (Google), Coq



RL for Normalization and Synthesis Tasks

- · Gauthier'19, 20: synthesizing simple programs and conjectures in logic
- Tree Neural Nets and RL (MCTS, policy/value) for:
- Guiding normalization in Robinson arithmetic
- · Guiding synthesis of combinators for a given lambda expression
- · Guiding synthesis of a diophantine equation characterizing a given set
- Quite encouraging results with a good curriculum (LPAR'20, CICM'20)
- · Motivated by TacticToe: argument synthesis/conjecturing is important
- · The results are series of applications of correct/explainable rules
- · Gauthier's deep RL framework verifies the whole series (proof) in HOL4

More on Synthesis and Conjecturing in Mathematics

- Targeted: generate intermediate lemmas (cuts) for a harder conjecture
- Unrestricted (theory exploration):
- · Creation of interesting conjectures based on the previous theory
- · One of the most interesting activities mathematicians do (how?)
- · Higher-level Al/reasoning task can we learn it?
- If so, we have solved math:
- ... just (recursively) divide Fermat into many subtasks ...
- ... and conquer (I mean: hammer) them away
- · Goes back to Langley (Bacon), Lenat (AM), Fajtlowicz (Graffiti)
- Combined with TP by Colton et al. in early 2000s (HR)
- Statistical methods, RNNs and Transformers by our groups since 2014

Can you find the flaw(s) in this fake GPT-2 proof?

| 🛛 Applications Places 🌍 | 🚞 🙆 🐏 4,71 GHz 🖇 | Wed 15:02 | Wed 15:02 |
|---|------------------|-----------|-----------|
| emacs@dell | | | • • • |
| File Edit Options Buffers Tools Index Mizar Hide/Show Help | | | _ |
| | | | |
| generated theorem with "proot" | | | |
| theorem Th23: :: STIRL2_1:23 | | | |
| for X, Y being finite set st not X is empty $\&$ X c= Y | | | |
| & card X = card Y holds X = Y | | | |
| proof | | | |
| let X, Y be finite set ; | | | |
| :: thesis: not X is empty & X c= Y & card X = card Y impli | es X = Y | | |
| assume that | | | |
| A1: not X is empty and A2: X c= Y and A3: card X = card | ίΥ; | | |
| :: thesis: $X = Y$ | | | |
| card $(Y \setminus X) = (card Y) - (card X)$ by A1, A3, CARD 2:44; | | | |
| then A4: card $(Y \setminus X) = ((card Y) - 1) - (card X)$ by CARD | 1:30; | | |
| $X = Y \setminus X$ by A2 A3 Th22. | | | |
| hence $X = Y$ by A4 XBOOLE 0 def 10 | | | |
| u thesis verum | | | |
| | | | |
| lena; | | | |
| -: card tst.miz 99% L2131 (Mizar Errors:13 hs Und | o-Tree) | | |

Figure: Fake full declarative GPT-2 "proof" - typechecks!

Proving the conditioned completions - MizAR hammer

| Applications Places | emacs@dell | Wed 14:42 | Wed 14:4 |
|--|--|-----------|----------|
| File Edit Options Buffers Tools Index Mizar Hide/Show Help P 🍙 🚍 🛪 🛤 Save 🔶 Undo 🔀 🌆 🛍 🔍 | | | |
| begin | | | |
| for M, N being Cardinal holds card M c= M V N by XBOOLE_1 | :7,CARD_3:44,CARD_1:7,CARD_1:3; :: [ATP details] | | |
| for X, Y being finite set st not X is empty & X c= Y & card X = | = card Y holds X = Y by CARD_FIN:1; :: [ATP details] | | |
| for M, N being Cardinal holds (M in N iff card M c= N) by Unsolved; :: [ATP details] | | | |
| for M, N being Cardinal holds (M in N iff card M in N) by CARD_3:44,CARD_1:9; :: [ATP det | tails] | | |
| for M, N being Cardinal holds Sum (M> N) = M $*$ N by CAF | RD_2:65; :: [ATP details] | | |
| for M, N being Cardinal holds M \wedge (union N) in N by Unsolved | d; :: [ATP details] | | |
| for M, N being Cardinal holds M *' N = N *' M by ATP-Unsolv | ved; :: [ATP details] | | |
| | | | |
| | | | |
| | | | |
| -: card_tst.miz 3% L47 (Mizar Errors:2 hs Undo-Tree |) | | |

Wrote /home/urban/mizwrk/7.13.01 4.181.1147/tst8/card tst.miz

Some GPT-2 conjectures

· Kinyon and Stanovsky (algebraists) confirmed that this cut is valid:

theorem Th10: :: GROUPP_1:10 for G being finite Group for N being normal Subgroup of G st N is Subgroup of center G & G ./. N is cyclic holds G is commutative

The generalization that avoids finiteness:

for G being Group for N being normal Subgroup of G st N is Subgroup of center G & G ./. N is cyclic holds G is commutative

· Some are clearly false, yet quite natural to ask:

```
theorem :: SINCOS10:17
sec is increasing on [0, pi/2)
leads to conjecturing the following:
Every differentiable function is increasing.
```

Gibberish Generator Provoking Algebraists

| Applic | ations Place | es 🌍 | | Group | o conjecture - | josef.urt | ban@gmail | l.com - Gmai | il - Chromiu | m | | - | <mark>0</mark> et 2 | ,28G | Hz 🎙 | Wed 1 | 7:12 | Wed 1 | 7:12 • • • |
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Figure: First successes in making mathematicians comment on AI.

Neural Autoformalization (Wang et al., 2018)



- · generate ca 1M Latex/Mizar pairs based on Bancerek's work
- train neural seq-to-seq translation models (Luong NMT)
- evaluate on about 100k examples, achieves 48% on unseen examples
- · many architectures tested, some work much better than others
- · very important latest invention: attention in the seq-to-seq models
- more data very important for neural training our biggest bottleneck
- Recent addition: unsupervised methods (Lample et all 2018) no need for aligned data!

| Rendered ⊮T _E X | Suppose s_8 is convergent and s_7 is convergent . Then $\lim(s_8+s_7) = \lim s_8 + \lim s_7$ |
|-------------------------------|---|
| Input LATEX | <pre>Suppose \$ { s _ { 8 } } \$ is convergent and \$ { s _ { 7 } } \$ is convergent . Then \$ \mathop { \rm lim } ({ s _ { 8 } } { + } { s _ { 7 } }) \mathrel { = } \mathop { \rm lim } { s _ { 8 } } { + } \mathop { \rm lim } { s _ { 7 } } \$.</pre> |
| Correct | <pre>seq1 is convergent & seq2 is convergent implies lim (seq1 + seq2) = (lim seq1) + (lim seq2) ;</pre> |
| Snapshot- 1000 | x in dom f implies (x * y) * (f (x (y (y y)))) = (x (y (y (y y))))) ; |
| Snapshot- 2000 | seq is summable implies seq is summable ; |
| Snapshot- 3000 | seq is convergent & lim seq = 0c implies seq = seq ; |
| Snapshot- 4000 | <pre>seq is convergent & lim seq = lim seq implies seq1 + seq2 is convergent ;</pre> |
| Snapshot- 5000 | <pre>seq1 is convergent & lim seq2 = lim seq2 implies lim_inf seq1 = lim_inf seq2 ;</pre> |
| Snapshot- 6000 | <pre>seq is convergent & lim seq = lim seq implies seq1 + seq2 is convergent ;</pre> |
| Snapshot- 7000 | seq is convergent & seq9 is convergent implies lim (seq + seq9) = (lim seq) + (lim seq9) ; |

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- Will be hybrid in 2021 as in 2020